

## Object orientation

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### Main components of object-orientation?

- Encapsulation
  - Group data and associated operations
- Information hiding
  - Need to know principle!
- Inheritance

## Inheritance

- `TYPE Shape = OBJECT`  
    `color: Colors;`  
    `END;`
- `TYPE Circle = Shape OBJECT`  
    `radius: INTEGER;`  
    `END;`

## One way to think about inheritance

- An object contains
  - All the fields of its supertype, followed by all the fields declared in the object
  - (All the methods of the supertype followed by all the methods in the object)
- What's a subtype?

## What do they look like in memory?

- Run time representation needs
  - space for data
  - (mapping of methods)
- What do Shape and Circle objects look like?
- How are they different from records?

## Another example

- **TYPE FilledCircle = Circle OBJECT**  
**fill\_color: Colors;**  
**END;**
- A FilledCircle is an object with
  - all the fields of a shape, followed by
  - all the fields of a circle, followed by
  - all the fields declared in FilledCircle
- FilledCircle <: Circle <: Shape

## Inheritance and subtyping

- Is a FilledCircle bigger or smaller than a Circle?
- Does this make sense from a subtyping point of view?

## An example

- **T = OBJECT ... END;**  
**S = T OBJECT ... END;**  
**x: REFANY;**  
**x := NEW(S);**  
**NARROW (x, T)**
- What representation does this need?

## Another example

- `l: ShapeList.T;`  
`l := ShapeList.Cons(NEW(Circle), l)`  
`l := ShapeList.Cons(NEW(FilledCircle), l)`  
`l := ShapeList.Cons(NEW(Square), l)`
- Is this legal?

## Continuing example

- `WHILE l.tail # NIL DO`  
`l.head.radius := 10`  
`l := l.tail;`  
`END`
- `WHILE l.tail # NIL DO`  
`IF l.head.color = Colors.White THEN`  
`...`  
`END;`  
`l := l.tail;`  
`END;`

## Continuing example

```
• PROCEDURE f(s: Shape) =  
  IF s.color # Colors.White THEN  
    ...  
  END
```

## Polymorphism

- Contrast with **monomorphic**:
  - Functions and their operands have a unique type
  - Every value and variable can be interpreted to be of one and only one type
- More directly
  - Functions work uniformly on a range of operand types
  - Some values and variables may have more than one type

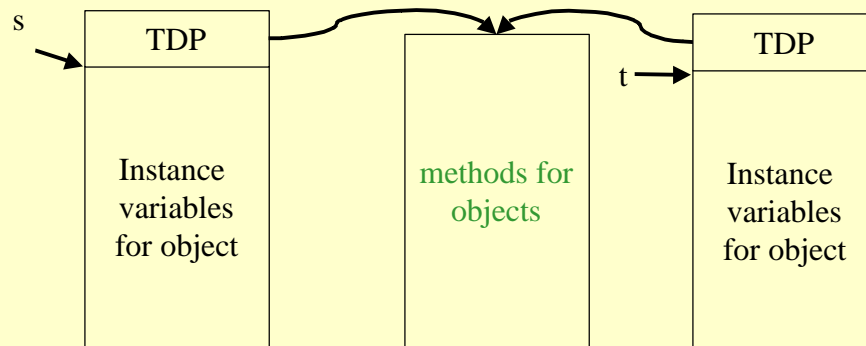
## Polymorphism versus monomorphism

- Expressiveness
- Simplicity
- Ease of implementation
- Efficiency

## Adding methods to objects: support for encapsulation

- `TYPE Shape = OBJECT`  
    `color: Colors;`  
    METHODS  
        `setColor(to: Color) := ...`  
    END;
- `TYPE Circle = Shape OBJECT`  
    `radius: INTEGER;`  
    END;

## Representation



## Let's look at the guts of a method

- `TYPE Shape = OBJECT`  
    `color: Colors;`  
    METHODS  
        `setColor(to: Color) := setColorShape;`  
    END;
- `PROCEDURE setColorShape(to: Color) =`  
    ...
- How does `setColorShape` access the "color" instance variable?

## Two approaches

- Java, C++, Smalltalk:
  - pass implicit parameter called “this” or “self”
- Modula-3:
  - s.setColor(c) => setColorShape(s, c)

## Overriding methods

- Overriding gives a new implementation for an existing method
- ```
TYPE Shape = OBJECT
  color: Colors;
METHODS render(); END;
TYPE Circle = Shape OBJECT
  radius: INTEGER;
OVERRIDES render := ...; END;
TYPE Square = Shape OBJECT
  side: INTEGER;
OVERRIDES render := ...; END;
```

## Representation

- Each type has its own mapping of methods to implementations
- The correct implementation is picked at run time!

## Some implications of inheritance

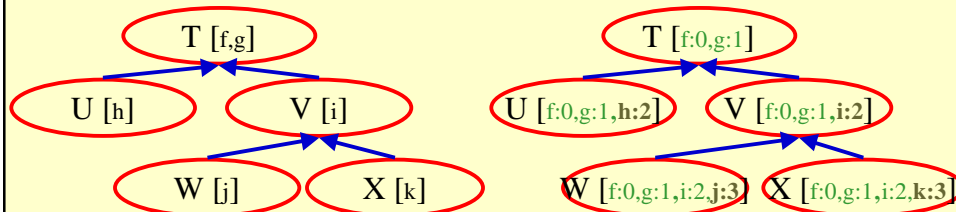
- ```
WHILE l.tail # NIL DO
  f(l.head)
  l := l.tail;
END;
```
- ```
PROCEDURE f(s: Shape) =
  s.render();
```
- What method is executed here?

## How to implement method dispatch?

- Assumptions:
  - Static typing
  - Single inheritance (we will look at multiple inheritance later)

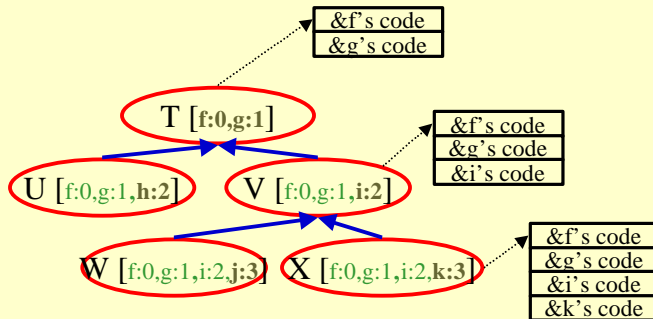
## V-Tables

- Idea:
  - Prepend the methods of a supertype to a subtype
  - A method **T::m** appears in the same position in all T's subclasses



## V-tables (cont.)

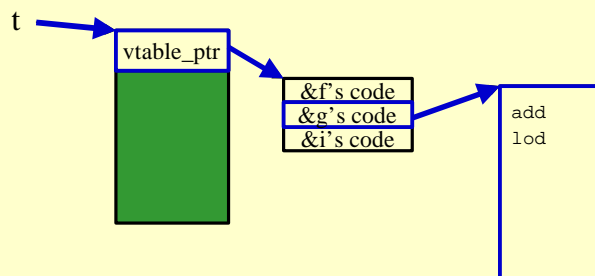
Construct a v-table for each class



v-tables are typically part of the type descriptor

## V-tables(cont.)

`t->g()` becomes  
`vp = t->vtable_ptr`  
`gaddr = *(vp+g's offset)`  
`(*gaddr)()`



## More examples

- `t: Shape;`  
`t := NEW Circle;`  
`t.render();`
- `t: Circle;`  
`t := NEW Circle;`  
`t.render();`  
`NARROW(t, Shape).render();`

## Overriding versus redefining

- `TYPE Shape = OBJECT`  
`color: Colors;`  
`METHODS render(); END;`  
`TYPE Circle = Shape OBJECT`  
`radius: INTEGER;`  
`METHODS render := ...; END;`  
`TYPE Square = Shape OBJECT`  
`side: INTEGER;`  
`METHODS render := ...; END;`

## Back to our old examples

- `WHILE l.tail # NIL DO`  
    `f(l.head)`  
`END;`
- `PROCEDURE f(s: Shape) =`  
    `s.render();`

## Back to our old examples (cont.)

- `t: Shape;`  
    `t := NEW Circle;`  
    `t.render();`
- `t: Circle;`  
    `t := NEW Circle;`  
    `t.render();`  
    `NARROW(t, Shape).render();`

## Contravariance and inheritance

- Issues with
  - The “id” function
  - The “clone” method

## Next topic: Case study

- Reading: Read up about objects in Java in your favorite Java book